

# HEDEN'S BOUND ON THE TAIL OF A VECTOR SPACE PARTITION

SASCHA KURZ\*

ABSTRACT. A vector space partition of  $\mathbb{F}_q^v$  is a collection of subspaces such that every non-zero vector is contained in a unique element. We improve a lower bound of Heden, in a subcase, on the number of elements of the smallest occurring dimension in a vector space partition. To this end, we introduce the notion of  $q^r$ -divisible sets of  $k$ -subspaces in  $\mathbb{F}_q^v$ . By geometric arguments we obtain non-existence results for these objects, which then imply the improved result of Heden.

## 1. INTRODUCTION

Let  $q > 1$  be a prime power,  $\mathbb{F}_q$  be the finite field with  $q$  elements, and  $v$  a positive integer. A *vector space partition*  $\mathcal{P}$  of  $\mathbb{F}_q^v$  is a collection of subspaces with the property that every non-zero vector is contained in a unique member of  $\mathcal{P}$ . If  $\mathcal{P}$  contains  $m_d$  subspaces of dimension  $d$ , then  $\mathcal{P}$  is of type  $k^{m_k} \dots 1^{m_1}$ . We may leave out some of the cases with  $m_d = 0$ . Subspaces of dimension  $d$  are also called *d-subspaces*. 1-subspaces are called *points*,  $(v-1)$ -subspaces are called *hyperplanes*, and each  $k$ -subspace contains  $\begin{bmatrix} k \\ 1 \end{bmatrix}_q := \frac{q^k-1}{q-1}$  points. So, in a vector space partition  $\mathcal{P}$  each point of the ambient space  $\mathbb{F}_q^v$  is covered by exactly one point of one of the elements of  $\mathcal{P}$ . An example of a vector space partition is given by a *k-spread* in  $\mathbb{F}_q^v$ , where  $\begin{bmatrix} v \\ 1 \end{bmatrix}_q / \begin{bmatrix} k \\ 1 \end{bmatrix}_q$   $k$ -subspaces partition the set of points of  $\mathbb{F}_q^v$ . The corresponding type is given by  $k^{m_k}$ , where  $m_k = \begin{bmatrix} v \\ 1 \end{bmatrix}_q / \begin{bmatrix} k \\ 1 \end{bmatrix}_q$ . If  $k$  divides  $v$  then considering the points of  $\mathbb{F}_q^{v/k}$  as  $k$ -dimensional subspaces over  $\mathbb{F}_q$  gives a construction of  $k$ -spreads. If  $k$  does not divide  $v$ , then no  $k$ -spreads exist. Vector space partitions of type  $k^{m_k} 1^{m_1}$  are known under the name *partial k-spreads*. More precisely, a partial  $k$ -spread in  $\mathbb{F}_q^v$  is a set  $\mathcal{K}$  of  $k$ -subspaces such that each point of the ambient space  $\mathbb{F}_q^v$  is covered at most by one of its elements. Adding the set of uncovered points, which are also called *holes*, gives a vector space partition of type  $k^{m_k} 1^{m_1}$ . Maximizing  $m_k = \#\mathcal{K}$  is equivalent to the minimization of  $m_1$ . If  $d_1$  is the smallest dimension with  $m_{d_1} \neq 0$ , we call  $m_{d_1}$  the *length of the tail* and call the set of the corresponding  $d_1$ -subspace the *tail*. Vector space partitions with a tail of small length are of special interest. In [4] Olof Heden obtained:

**Theorem 1.** (Theorem 1 in [4]) Let  $\mathcal{P}$  be a vector space partition of type  $d_1^{u_1} \dots d_2^{u_2} d_1^{u_1}$  of  $\mathbb{F}_q^v$ , where  $u_1, u_2 > 0$  and  $d_1 > \dots > d_2 > d_1 \geq 1$ .

- (i) If  $q^{d_2-d_1}$  does not divide  $u_1$  and if  $d_2 < 2d_1$ , then  $u_1 \geq q^{d_1} + 1$ ;
- (ii) if  $q^{d_2-d_1}$  does not divide  $u_1$  and if  $d_2 \geq 2d_1$ , then either  $d_1$  divides  $d_2$  and  $u_1 = \begin{bmatrix} d_2 \\ 1 \end{bmatrix}_q / \begin{bmatrix} d_1 \\ 1 \end{bmatrix}_q$  or  $u_1 > 2q^{d_2-d_1}$ ;
- (iii) if  $q^{d_2-d_1}$  divides  $u_1$  and  $d_2 < 2d_1$ , then  $u_1 \geq q^{d_2} - q^{d_1} + q^{d_2-d_1}$ ;
- (iv) if  $q^{d_2-d_1}$  divides  $u_1$  and  $d_2 \geq 2d_1$ , then  $u_1 \geq q^{d_2}$ .

Moreover, in Theorem 2 and Theorem 3 he classified the possible sets of  $d_1$ -subspaces for  $u_1 = q^{d_1} + 1$  and  $u_1 = \begin{bmatrix} d_2 \\ 1 \end{bmatrix}_q / \begin{bmatrix} d_1 \\ 1 \end{bmatrix}_q$ , respectively. The results were obtained using the theory of mixed perfect 1-codes, see e.g. [6].

In [2] the authors improved a lower bound of Heden [2] on the size of inclusion-maximal partial 2-spreads by translating the underlying techniques into geometry. Here we improve Theorem 1(ii). The underlying geometric structure is the set  $\mathcal{N}$  of  $d_1$ -subspaces of a vector space partition  $\mathcal{P}$  of type  $d_1^{u_1} \dots d_2^{u_2} d_1^{u_1}$ . For  $d_1$  this is just a set of points in  $\mathbb{F}_q^v$ . It can be shown that the existence of  $\mathcal{P}$  implies  $\#\mathcal{N} \equiv \#(\mathcal{N} \cap H) \pmod{q^{d_2-1}}$  for every hyperplane  $H$  of  $\mathbb{F}_q^v$ , see e.g. [7]. Taking a vector representation of the elements of  $\mathcal{N}$  as columns of a generator matrix, we obtain a corresponding (projective) linear code  $\mathcal{C}$  over  $\mathbb{F}_q$ . The modulo constraints for  $\mathcal{N}$  are equivalent to the property that the Hamming weights of the codewords of  $\mathcal{C}$  are divisible by  $q^{d_2-1}$ . The study of so-called divisible codes, where the Hamming weights of the codewords of a linear code are divisible by some factor  $\Delta > 1$ , was initiated by Harold Ward, see e.g. [9]. The MacWilliams identities, linking the weight distribution of a linear code with the weight distribution of its dual code, can be relaxed to a linear program. Incorporating some information about the weight distribution of a linear code may result in an infeasible linear program, which then certifies the non-existence of such a code. This technique is known under the name linear programming method for codes and was more generally developed for association schemes by Philip Delsarte [3]. In [8] analytic solutions of linear programs for projective  $q^r$ -divisible linear codes have been applied in order to compute upper bounds for partial  $k$ -spreads. Indeed, all currently known upper bounds for partial  $k$ -spreads can be deduced from this method, see [7] for a survey.

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Here, we generalize the approach to the case  $d_1 > 1$  by studying the properties of the set  $\mathcal{N}$  of  $d_1$ -subspaces of a vector space partition  $\mathcal{P}$  of  $\mathbb{F}_q^v$  of type  $d_1^{u_1} \dots d_2^{u_2} d_1^{u_1}$  in Section 2. It turns out that we have  $\#\mathcal{N} \equiv \#(\mathcal{N} \cap H) \pmod{q^{d_2-d_1}}$  for every hyperplane  $H$  of  $\mathbb{F}_q^v$ , see Lemma 3, which we introduce as a definition of a  $q^{d_2-d_1}$ -divisible set of  $k$ -subspaces with trivial intersection. By elementary counting techniques we obtain a partial substitute for the MacWilliams identities, see the equations (1) and (2). These imply some analytical criteria for the non-existence of such sets  $\mathcal{N}$ , which are used in Section 3 to reprove Theorem 1. By an improved analysis we tighten Theorem 1 to Theorem 12. More precisely, the second lower bound of Theorem 1(ii) is improved. We close with some numerical results on the spectrum of the possible cardinalities of  $\mathcal{N}$  and pose some open problems.

## 2. SETS OF DISJOINT $k$ -SUBSPACES AND THEIR INCIDENCES WITH HYPERPLANES

For a positive integer  $k$  let  $\mathcal{N}$  be a set of pairwise disjoint, i.e., having trivial intersection,  $k$ -subspaces in  $\mathbb{F}_q^v$ , where we assume that the  $k$ -subspaces from  $\mathcal{N}$  span  $\mathbb{F}_q^v$ , i.e.,  $v$  is minimally chosen. By  $a_i$  we denote the number of hyperplanes  $H$  of  $\mathbb{F}_q^v$  with  $\#(\mathcal{N} \cap H) := \#\{U \in \mathcal{N} : U \leq H\} = i$  and set  $n := \#\mathcal{N}$ . Due to our assumption on the minimality of the dimension  $v$  not all  $n$  elements from  $\mathcal{N}$  can be contained in a hyperplane. Double-counting the incidences of the tuples  $(H)$ ,  $(B_1, H)$ , and  $(B_1, B_2, H)$ , where  $H$  is a hyperplane and  $B_1 \neq B_2$  are elements of  $\mathcal{N}$  contained in  $H$  gives:

$$\sum_{i=0}^{n-1} a_i = \begin{bmatrix} v \\ 1 \end{bmatrix}_q, \quad \sum_{i=0}^{n-1} i a_i = n \cdot \begin{bmatrix} v-k \\ 1 \end{bmatrix}_q, \quad \text{and} \quad \sum_{i=0}^{n-1} i(i-1) a_i = n(n-1) \cdot \begin{bmatrix} v-2k \\ 1 \end{bmatrix}_q. \quad (1)$$

For three different elements  $B_1, B_2, B_3$  of  $\mathcal{N}$  their span  $\langle B_1, B_2, B_3 \rangle$  has a dimension  $i$  between  $2k$  and  $3k$ . Denoting the number of corresponding triples by  $b_i$ , double-counting tuples  $(B_1, B_2, B_3, H)$ , where  $H$  is a hyperplane and  $B_1, B_2, B_3$  are pairwise different elements of  $\mathcal{N}$  contained in  $H$ , gives:

$$\sum_{i=0}^{n-1} i(i-1)(i-2) a_i = \sum_{i=2k}^{3k} b_i \begin{bmatrix} v-i \\ 1 \end{bmatrix}_q \quad \text{and} \quad \sum_{i=2k}^{3k} b_i = n(n-1)(n-2). \quad (2)$$

Given parameters  $q, k, n$ , and  $v$  the so-called (*integer*) *linear programming method* asks for a solution of the equation system given by (1) and (2) with  $a_i, b_i \in \mathbb{R}_{\geq 0}$  ( $a_i, b_i \in \mathbb{N}$ ). If no solution exists, then no corresponding set  $\mathcal{N}$  can exist. For  $k = 1$  the equations from (1) and (2) correspond to the first four MacWilliams identities, see e.g. [7].

If there is a single non-zero value  $a_i$  the system can be solved analytically.

**Lemma 2.** *If  $a_i = 0$  for all  $i \neq r > 0$  and  $k < v$  in the above setting, then there exists an integer  $s \geq 2$  with  $v = sk$  and  $\mathcal{N}$  is a  $k$ -spread. Additionally we have  $r = \frac{q^{v-k}-1}{q^k-1}$ .*

PROOF. Solving (1) for  $r, a_r$ , and  $n$  gives  $n = \frac{q^{2v-k} - q^v - q^{v-k} + 1}{q^v - q^{v-k} - q^k + 1}$ . Writing  $v = sk + t$  with  $s, t \in \mathbb{N}$  and  $0 \leq t < k$  we obtain  $n = \sum_{i=1}^s q^{v-ik} + \frac{q^{v-k+t} - q^{v-k} - q^t + 1}{q^v - q^{v-k} - q^k + 1}$ . Since  $n \in \mathbb{N}$  and  $0 \leq q^{v-k+t} - q^{v-k} - q^t + 1 < q^v - q^{v-k} - q^k + 1$  we have  $q^{v-k+t} - q^{v-k} - q^t + 1 = 0$  so that  $t = 0$  and  $n = \frac{q^v-1}{q^k-1}$ . Counting points gives that  $\mathcal{N}$  partitions  $\mathbb{F}_q^v$ .  $\square$

We remark that  $r = 0$  forces  $n \in \{0, 1\}$  so that  $\mathcal{N}$  is empty or consists of a single  $k$ -subspace in  $\mathbb{F}_q^k$  and  $v = k$  implies the latter case. So, these degenerated cases correspond to  $s \in \{0, 1\}$  in Lemma 2. As pointed out after [4, Theorem 2], such results can be proved in different ways. While the case that only one  $a_i$  is non-zero is rather special, we can show that many  $a_i$  are equal to zero in our setting.

**Lemma 3.** *Let  $\mathcal{P}$  be a vector space partition of type  $d_1^{u_1} \dots d_2^{u_2} d_1^{u_1}$  of  $\mathbb{F}_q^v$ , where  $u_1, u_2 > 0$ , and let  $\mathcal{N}$  be the set of  $d_1$ -subspaces. Then, we have  $\#\mathcal{N} \equiv \#(\mathcal{N} \cap H) \pmod{q^{d_2-d_1}}$  for every hyperplane  $H$  of  $\mathbb{F}_q^v$ .*

PROOF. For each  $U \in \mathcal{P}$  we have  $\dim(U \cap H) \in \{\dim(U), \dim(U) - 1\}$ . So counting points in  $\mathbb{F}_q^v$  and  $H$  gives the existence of integers  $a, a'$  with  $m \cdot \begin{bmatrix} d_2 \\ 1 \end{bmatrix}_q + a q^{d_2} + u_1 \begin{bmatrix} d_1 \\ 1 \end{bmatrix}_q = \begin{bmatrix} v \\ 1 \end{bmatrix}_q$  and  $m \cdot \begin{bmatrix} d_2-1 \\ 1 \end{bmatrix}_q + a' q^{d_2-1} + u_1' q^{d_1-1} + u_1 \begin{bmatrix} d_1-1 \\ 1 \end{bmatrix}_q = \begin{bmatrix} v-1 \\ 1 \end{bmatrix}_q$ , where  $m := \sum_{i=2}^l u_i$  and  $u_1' := \#(\mathcal{N} \cap H)$ . By subtraction we obtain  $m q^{d_2-1} + a q^{d_2} - a' q^{d_2-1} + u_1 q^{d_1-1} - u_1' q^{d_1-1} = q^{v-1}$ , so that  $u_1 q^{d_1-1} \equiv u_1' q^{d_1-1} \pmod{q^{d_2-1}}$ .  $\square$

**Definition 4.** Let  $\mathcal{N}$  be a set of  $k$ -subspaces in  $\mathbb{F}_q^v$ . If there exists a positive integer  $r$  such that  $a_i$  is non-zero only if  $\#\mathcal{N} - i$  is divisible by  $q^r$  and the  $k$ -subspaces are pairwise disjoint, then we call  $\mathcal{N}$   $q^r$ -divisible.

Using the notation of Lemma 3,  $\mathcal{N}$  is  $q^{d_2-d_1}$ -divisible. As mentioned in the introduction, for  $d_1 = 1$ , taking the elements of  $\mathcal{N}$  as columns of a generator matrix, we obtain a projective linear code, whose Hamming weights are divisible by  $q^{d_2-1}$ .

**Example 5.** For integers  $k \geq 2$  and  $r = ak + b$  with  $0 \leq b < k$  let  $\mathcal{N}$  be a  $k$ -spread of  $\mathbb{F}_q^{(a+2)k}$ . Starting from a  $(a+2)k$ -spread in  $\mathbb{F}_q^{2(a+2)k}$  we obtain a vector space partition  $\mathcal{P}$  by replacing one  $(a+2)k$ -dimensional spread

element with  $\mathcal{N}$ . From Lemma 3 and  $q^r | q^{(a+2)k-k} = q^{(a+1)k}$  we deduce that the set  $\mathcal{N}$  of  $k$ -subspaces is  $q^r$ -divisible. Its cardinality is given by  $\binom{(a+2)k}{1}_q / \binom{k}{1}_q$ .

**Example 6.** For integers  $k \geq 2$  and  $r \geq 1$  let  $n = k + r$  and consider a matrix representation  $M: \mathbb{F}_{q^n} \rightarrow \mathbb{F}_q^{n \times n}$  of  $\mathbb{F}_{q^n}/\mathbb{F}_q$ , obtained by expressing the multiplication maps  $\mu_\alpha: \mathbb{F}_{q^n} \rightarrow \mathbb{F}_{q^n}$ ,  $x \mapsto \alpha x$ , which are linear over  $\mathbb{F}_q$ , in terms of a fixed basis of  $\mathbb{F}_{q^n}/\mathbb{F}_q$ . Then, all matrices in  $M(\mathbb{F}_{q^n})$  are invertible and have mutual rank distance  $d_R(A, B) := \text{rk}(A - B) = n$ , see e.g. [7] for proofs of these and the subsequent facts. In other words, the matrices of  $M(\mathbb{F}_{q^n})$  form a maximum rank distance code with minimum rank distance  $n$  and cardinality  $q^n$ .

Now let  $\mathcal{B} \subseteq \mathbb{F}_q^{k \times n}$  be the matrix code obtained from  $M(\mathbb{F}_{q^n})$  by deleting the last  $n - k$  rows, say, of every matrix. Then  $\mathcal{B}$  has cardinality minimum rank distance  $k$ . Hence, by applying the lifting construction  $B \mapsto (I_k | B)$ , where  $I_k$  is the  $k \times k$  identity matrix, to  $\mathcal{B}$  we obtain a partial  $k$ -spread  $\mathcal{N}$  in  $\mathbb{F}_q^v$  of size  $q^n = q^{k+r}$ . Since precisely the points outside the  $(k+r)$ -subspace  $S = \{x \in \mathbb{F}_q^v : x_1 = x_2 = \dots = x_k = 0\}$  are covered,  $\mathcal{P} = \mathcal{N} \cup \{S\}$  is a vector space partition of  $\mathbb{F}_q^{2k+r}$  and  $\mathcal{N}$  is  $q^{k+r}$ -divisible with cardinality  $q^{k+r}$ .

From the first two equations of (1) we deduce:

**Lemma 7.** For a  $q^r$ -divisible set  $\mathcal{N}$  of  $k$ -subspaces in  $\mathbb{F}_q^v$ , there exists a hyperplane  $H$  with  $\#(\mathcal{N} \cap H) \leq n/q^k$ .

PROOF. Let  $i$  be the smallest index with  $a_i \neq 0$ . Then, the first two equations of (1) are equivalent to  $\sum_{j \geq 0} a_{i+q^r j} = \binom{v}{1}_q$  and  $\sum_{j \geq 0} (i + q^r j) \cdot a_{i+q^r j} = n \binom{v-k}{1}_q$ . Subtracting  $i$  times the first equation from the second equation gives  $\sum_{j > 0} q^r j a_{i+q^r j} = n \cdot \frac{q^{v-k}-1}{q-1} - i \cdot \frac{q^v-1}{q-1}$ . Since the left-hand side is non-negative, we have  $i \leq \frac{q^{v-k}-1}{q^{v-1}-1} \cdot n \leq \frac{n}{q^k}$ .  $\square$

Stated less technical, the proof of Lemma 7 is given by the fact that the hyperplane with the minimum number of  $k$ -subspaces contains at most as many  $k$ -subspaces as the average number of  $k$ -subspaces per hyperplane.

Taking also the third equation of (1) into account implies a quadratic criterion:

**Lemma 8.** Let  $m \in \mathbb{Z}$  and  $\mathcal{N}$  be a  $q^r$ -divisible set of  $k$ -subspaces in  $\mathbb{F}_q^v$ . Then,  $\tau(n, q^r, q^k, m) \cdot q^{v-2k-2r} - m(m-1) \geq 0$ , where  $\tau(n, \Delta, u, m) := \Delta^2 u^2 m(m-1) - n(2m-1)u(u-1)\Delta + n(u-1)(n(u-1)+1)$ .

PROOF. With  $y = q^{v-2k}$ ,  $u = q^k$ , and  $\Delta = q^r$ , we can rewrite the equations of (1) to  $u^2 y - 1 = (q-1) \sum_{i \in \mathbb{Z}} a_i$ ,  $n \cdot (uy - 1) = (q-1) \sum_{i \in \mathbb{Z}} i a_i$ , and  $n(n-1) \cdot (y-1) = \sum_{i \in \mathbb{Z}} i(i-1) a_i$ .  $(n-m\Delta)(n-(m-1)\Delta)$  times the first minus  $2n - (2m-1)\Delta - 1$  times the second plus the third equation gives  $y \cdot \tau(n, \Delta, u, m) - \Delta^2 m(m-1) = (q-1) \sum_{i \in \mathbb{Z}} (n-m\Delta-i)(n-(m-1)\Delta-i) a_i = (q-1) \sum_{h \in \mathbb{Z}} \Delta^2 (m-h)(m-h+1) a_{n-h\Delta} \geq 0$ .  $\square$

As a preparation we present another classification result:

**Lemma 9.** If  $\mathcal{N}$  is a  $q$ -divisible set of  $k$ -subspaces in  $\mathbb{F}_q^v$  of cardinality  $q^k + 1$ , then  $\mathcal{N}$  partitions  $\mathbb{F}_q^{2k}$ .

PROOF. Setting  $c_i := (q-1)a_{1+iq}$  and  $l := q^{k-1} - 1$  we can rewrite the equations of (1) to  $\sum_{i=0}^l c_i = q^v - 1$ ,  $\sum_{i=0}^l (1+iq)c_i = (q^k+1)(q^{v-k}-1)$ , and  $\sum_{i=0}^l iq(1+iq)c_i = (q^k+1)q^k(q^{v-2k}-1)$ . Since  $ql+1$  times the second minus  $ql+1$  times the first minus the third equation gives  $0 \leq \sum_{i=0}^l iq^2(l-i)c_i = -q^{k+1}(q^{v-2k}-1)$ , we have  $v = 2k$ . Every point of  $\mathbb{F}_q^v$  is covered by an element from  $\mathcal{N}$  due to  $\binom{2k}{1}_q / \binom{k}{1}_q = q^k + 1$ .  $\square$

### 3. PROOF OF HEDEN'S RESULTS AND FURTHER IMPROVEMENTS

Let  $\mathcal{P}$  be a vector space partition of type  $d_1^{u_1} \dots d_2^{u_2} d_1^{u_1}$  of  $\mathbb{F}_q^v$ , where  $u_1, u_2 > 0$ ,  $d_1 > \dots > d_2 > d_1 \geq 1$ . Let  $\mathcal{N}$  be the set of  $d_1$ -subspaces and  $V$  be the subspace spanned by  $\mathcal{N}$ . By  $n$  we denote the cardinality of  $\mathcal{N}$  and by  $a_i$  we denote the number of hyperplanes of  $V$  that contain exactly  $i$  elements from  $\mathcal{N}$ .

Assume that  $q^{d_2-d_1}$  does not divide  $u_1$ . We have  $\#(\mathcal{N} \cap H) \geq 1$  for every hyperplane  $H$  of  $V$  due to Lemma 3, so that Lemma 7 gives  $u_1 \geq q^{d_1}$ . Thus, we have  $u_1 \geq q^{d_1} + 1$ . If  $u = q^{d_1} + 1$  then we can apply Lemma 9 for the classification of the possible sets  $\mathcal{N}$ . If  $u_1 < 2q^{d_2-d_1}$  then for  $a_i > 0$  we have  $i < q^{d_2-d_1}$  and  $i \equiv u_1 \pmod{q^{d_2-d_1}}$  so that we can apply Lemma 2. Thus, either  $d_2$  divides  $d_1$  and  $u_1 = (q^{d_2}-1)/(q^{d_1}-1)$  or  $u_1 > 2q^{d_2-d_1}$ . The first case can be attained by a  $d_2$ -spread where one  $d_2$ -subspace is replaced by a  $d_1$ -spread, see Example 5. We remark that no assumption on the relation between  $d_2$  and  $d_1$  is used in our derivation. However, if  $d_2 < 2d_1$  then  $d_1$  cannot divide  $d_2$  and  $q_1^{d_1} + 1 > 2q^{d_2-d_1}$ .

Assume that  $q^{d_2-d_1}$  divides  $u_1$ . Setting  $\Delta = q^{d_2-d_1}$ ,  $u = q^{d_1}$ ,  $n = \Delta l$ , and  $m = l^\dagger$  for some integer  $l$ , we conclude  $\tau(n, \Delta, u, m) = \Delta l(\Delta l - \Delta u + u - 1) \geq 0$  from Lemma 8, so that  $l \geq \lceil u - \frac{u}{\Delta} + \frac{1}{\Delta} \rceil$ . The right-hand side is equal to  $u = q^{d_1}$  if  $d_2 \geq 2d_1$  and to  $u - u/\Delta + 1 = q^{d_1} - q^{2d_1-d_2} + 1$  otherwise, which is equivalent to  $n \geq q^{d_2}$  and  $n \geq q^{d_2} - q^{d_1} + q^{d_2-d_1}$ . We remark that equality is achievable in the latter case via the 2-weight codes constructed in [1] (with parameters  $n' = d_1$  and  $m = d_2 - d_1$ ). We do not know whether the corresponding  $q^{d_2-d_1}$ -divisible set of  $d_1$ -subspaces can be realized as a vector space partition of  $\mathbb{F}_q^v$ .<sup>‡</sup> For the first case see Example 6.

<sup>†</sup>The choice for  $m$  can be obtained by minimizing  $\tau(n, \Delta, u, m)$ , i.e., solving  $\frac{\partial \tau(n, \Delta, u, m)}{\partial m} = 0$  and rounding.

<sup>‡</sup>A suitable test case might be to decide whether a vector space partition of type  $4^4 3^{135} 2^6$  exists in  $\mathbb{F}_2^{10}$ .

The above comprises [4, Theorems 1-4]. Given the stated examples, just Theorem 1(ii), for the case where  $d_1$  does not divide  $d_2$ , leaves some space for improving the lower bound on  $u_1$ . To that end we analyze Lemma 8 in more detail. Since the statements look rather technical and complicated we first give a justification for the necessity of this fact. Via the quadratic inequality of Lemma 8 intervals of cardinalities can be excluded for different values of the parameter  $m$ . However, some cardinalities are indeed feasible. If  $r = ak + b$  with  $0 \leq b < k$  then the two constructions from Example 5 and Example 6 give  $q^r$ -divisible set of  $k$ -subspaces of cardinality  $\begin{bmatrix} (a+2)k \\ 1 \end{bmatrix}_q / \begin{bmatrix} k \\ 1 \end{bmatrix}_q$  and  $q^{k+r}$ , respectively. For  $q = 2, r = 3, k = 2$  the cardinalities of these two examples are given by 21 and 32. In general, each two  $q^r$ -divisible sets  $\mathcal{N}_1$  and  $\mathcal{N}_2$  of  $k$ -subspaces can be combined to a  $q^r$ -divisible set of  $k$ -subspaces of cardinality  $\#\mathcal{N}_1 + \#\mathcal{N}_2$ . Since  $\begin{bmatrix} (a+2)k \\ 1 \end{bmatrix}_q / \begin{bmatrix} k \\ 1 \end{bmatrix}_q$  and  $q^{k+r}$  are coprime there exists some integer  $F_q(k, r)$  such that  $q^r$ -divisible sets of  $k$ -subspaces exist for every cardinality  $n > F_q(k, r)$ . Below that number some cardinalities can be excluded, but their density decreases with increasing  $n$ . Our numerical example is continued after the proof of Theorem 12.

**Proposition 10.** *Let  $\mathcal{N}$  be a  $q^r$ -divisible set of  $k$ -subspaces in  $\mathbb{F}_q^v$ ,  $u = q^k$  and  $\Delta = q^r$ . Then,  $n \notin \left[1, \frac{q^{k+r}-1}{q^r-1}\right)$  and*

$$n \notin \left[ \left[ \frac{1}{u-1} \cdot \left( \Delta um - \frac{\Delta u + 1}{2} - \frac{1}{2} \sqrt{\omega} \right) \right], \left[ \frac{1}{u-1} \cdot \left( \Delta um - \frac{\Delta u + 1}{2} + \frac{1}{2} \sqrt{\omega} \right) \right] \right],$$

where  $\omega = (\Delta u - 2m)^2 + (2\Delta u + 1 - 4m^2)$ , for all  $m \in \mathbb{N}$  with  $2 \leq m \leq \lfloor \frac{\Delta u}{4} + \frac{1}{2} + \frac{1}{4\Delta u} \rfloor$ .

PROOF. We set  $\bar{\Delta} = \Delta u$  and  $\bar{n} = n(u-1)$  so that  $\tau(n, \Delta, u, m) = \bar{\Delta}^2 m(m-1) - \bar{n}\bar{\Delta}(2m-1) + \bar{n}(\bar{n}+1)$ . We have  $\tau(n, \Delta, u, m) \leq 0$  iff  $\left| \bar{n} - \bar{\Delta}m + \frac{\bar{\Delta}+1}{2} \right| \leq \frac{1}{2} \sqrt{\bar{\Delta}^2 - 4m\bar{\Delta} + 2\bar{\Delta} + 1}$  and  $m \leq \frac{\bar{\Delta}}{4} + \frac{1}{2} + \frac{1}{4\bar{\Delta}}$ . Rewriting and applying Lemma 8 with  $1 \leq m \leq \lfloor \frac{\Delta u}{4} + \frac{1}{2} + \frac{1}{4\Delta u} \rfloor$  gives the result since  $m(m-1) > 0$  for  $m \geq 2$ .  $\square$

**Proposition 11.** *Let  $\mathcal{N}$  be a  $q^r$ -divisible set of  $k$ -subspaces in  $\mathbb{F}_q^v$ , where  $r = ak + b$  with  $a, b \in \mathbb{N}, 0 < b < k$  and  $a \geq 1$ . Then,  $n \geq \frac{q^{(a+2)k}-1}{q^k-1} = q^r \cdot q^{k-b} + \frac{q^r \cdot q^{k-b}-1}{q^k-1} = \Delta q^{k-b} + q^k \Theta + 1$ , where  $\Delta := q^r$  and  $\Theta := \frac{q^{ak}-1}{q^k-1}$ .*

PROOF. From Lemma 2 we conclude  $n \geq 2q^r$  and set  $u = q^k$ . For  $2 \leq m \leq q^{k-b}$  we have  $2\Delta u + 1 - 4m^2 > 0$ , so that Proposition 10 gives  $n \notin \left[ \left[ \frac{\Delta u(m-1)-1/2+m}{u-1} \right], \left[ \frac{\Delta um-1/2-m}{u-1} \right] \right]$ . Since  $\Delta(m-1) \leq \left\lceil \frac{\Delta u(m-1)-1/2+m}{u-1} \right\rceil = \Delta(m-1) + \left\lceil \frac{\Delta(m-1)-1/2+m}{u-1} \right\rceil \leq \Delta m$  and  $\left\lfloor \frac{\Delta um-1/2-m}{u-1} \right\rfloor = \Delta m + mq^b \Theta + \left\lfloor \frac{mq^b-1/2-m}{q^k-1} \right\rfloor = \Delta m + mq^b \Theta$ , we conclude  $n \notin [\Delta m, \Delta m + mq^b \Theta]$  for  $2 \leq m \leq q^{k-b}$ .

It remains to show  $n \notin [\Delta m, \Delta m + mq^b \Theta + 1, \Delta(m+1) - 1] =: I_m$  for all  $2 \leq m \leq q^{k-b} - 1$ . If  $n \in I_m$ , then we can write  $n = \Delta m + mq^b \Theta + x$  with  $x \geq 1$  and  $mq^b \Theta + x < \Delta$ , so that  $q^k \cdot (mq^b \Theta + x) = \Delta m + mq^b \Theta + (xq^k - mq^b) < \Delta m + mq^b \Theta + x = n$ , which contradicts Lemma 7.  $\square$

In other words, in the case of Theorem 1(ii), where  $d_2 = ad_1 + b$  with  $0 < b < d_1$  and  $a, b \in \mathbb{N}$ , we have  $u_1 \geq q^{d_2-d_1} \cdot q^{d_1-b} + \frac{q^{(a+1)d_1}-1}{q^{d_1}-1} = \frac{q^{(a+2)d_1}-1}{q^{d_1}-1}$ , which can be attained by an  $d_1$ -spread in  $\mathbb{F}_q^{(a+2)d_1}$ . Without the knowledge of  $b$ , we can state  $u_1 \geq q \cdot q^{d_2-d_1} + \left\lceil \frac{q^{d_2+1}-1}{q^{d_1}-1} \right\rceil$ , which also improves Theorem 1(ii) and is tight whenever  $d_2 + 1$  is divisible by  $d_1$ . Summarizing our findings we obtain our main theorem:

**Theorem 12.** *For a non-empty  $q^r$ -divisible set  $\mathcal{N}$  of  $k$ -subspaces in  $\mathbb{F}_q^v$  the following bounds on  $n = \#\mathcal{N}$  are tight.*

- (i) *We have  $n \geq q^k + 1$  and if  $r \geq k$  then either  $k$  divides  $r$  and  $n \geq \frac{q^{k+r}-1}{q^k-1}$  or  $n \geq \frac{q^{(a+2)k}-1}{q^k-1}$ , where  $r = ak + b$  with  $0 < b < k$  and  $a, b \in \mathbb{N}$ .*
- (ii) *Let  $q^r$  divide  $n$ . If  $r < k$  then  $n \geq q^{k+r} - q^k + q^r$  and  $n \geq q^{k+r}$  otherwise.*

For (i) the lower bounds are attained by  $k$ -spreads, see Example 5. For (ii) the second lower bound is attained by a construction based on lifted MRD codes, see Example 6. In the other case the 2-weight codes constructed in [1] attain the lower bound. Thus, Theorem 12 is tight and implies an improvement of Theorem 1(ii).

While the smallest cardinality of a non-empty  $q^r$ -divisible set of  $k$ -subspaces over  $\mathbb{F}_q$  has been determined, the spectrum of possible cardinalities remains widely unknown. For  $k = 1$  [7, Theorem 12] states that either  $n > rq^{r+1}$  or there exist integers  $a, b$  with  $n = a \begin{bmatrix} r+1 \\ 1 \end{bmatrix}_q + bq^{r+1}$  and bounds for the maximum excluded cardinality have been determined in [5]. However, Lemma 7 and Lemma 8, applied via Proposition 10, give restrictions going far beyond Theorem 12. For  $q = 2, r = 3, k = 2$ , and  $n \leq 81$  we exemplarily state that only  $n \in \{21, 31, 32, 33, 42, 43, 44, 52, \dots, 55, 62, \dots, 66, 72, \dots, 78\}$  might be attainable. The mentioned constructions cover the cases  $n \in \{21, 32, 42, 53, 63, 64, 74\} \subseteq \{21a + 32b : a, b \in \mathbb{N}\}$ . Replacing the lines by their contained 3 points, we obtain  $2^4$ -divisible sets of 1-subspaces in  $\mathbb{F}_q^v$  of cardinality  $3n$ , for which two further exclusion criteria have been presented in [7], excluding the cases  $n \in \{33, 44\}$ . [7, Lemma 23] is based on a cubic polynomial obtained from (1) and (2), similar to the quadratic polynomial from Lemma 8 obtained from (1). Here, the presence of  $k$  additional  $b_i$ -variables

may make the analysis more difficult for  $k > 1$ . For a  $q^r$ -divisible set  $\mathcal{N}$  of 1-subspaces we have that  $\mathcal{N} \cap H$  is  $q^{r-1}$ -divisible for every hyperplane  $H$ , which allows a recursive application of the linear programming method. For  $k > 1$  we need to consider  $k$ -subspaces and  $k - 1$ -subspaces in  $H$ , see [7, Section 6.3], which makes the bookkeeping more complicated.

The determination of the possible spectrum of cardinalities of  $q^r$ -divisible sets of  $k$ -subspaces remains an interesting open problem. Even for small parameters this might be challenging. A possible intermediate step is the determination of the number  $F_q(k, r)$  being similar to the Frobenius number. Extending the small list of constructions is also worthwhile.

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DEPARTMENT OF MATHEMATICS, UNIVERSITY OF BAYREUTH, 95440 BAYREUTH, GERMANY  
*E-mail address:* sascha.kurz@uni-bayreuth.de