

UPPER BOUNDS FOR PARTIAL SPREADS

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ABSTRACT. A *partial t -spread* in \mathbb{F}_q^n is a collection of t -dimensional subspaces with trivial intersection such that each non-zero vector is covered at most once. We present some improved upper bounds on the maximum sizes.

Keywords: Galois geometry, partial spreads, constant dimension codes, and vector space partitions

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1. INTRODUCTION

Let $q > 1$ be a prime power and n a positive integer. A *vector space partition* \mathcal{P} of \mathbb{F}_q^n is a collection of subspaces with the property that every non-zero vector is contained in a unique member of \mathcal{P} . If \mathcal{P} contains m_d subspaces of dimension d , then \mathcal{P} is of type $k^{m_k} \dots 1^{m_1}$. We may leave out some of the cases with $m_d = 0$. Subspaces of dimension 1 are called *holes*. If there is at least one non-hole, then \mathcal{P} is called non-trivial.

A *partial t -spread* in \mathbb{F}_q^n is a collection of t -dimensional subspaces such that the non-zero vectors are covered at most once, i.e., a vector space partition of type $t^{m_t} 1^{m_1}$. By $A_q(n, 2t; t)$ we denote the maximum value of m_t ¹. Writing $n = kt + r$, with $k, r \in \mathbb{N}_0$ and $r \leq t - 1$, we can state that for $r \leq 1$ or $n \leq 2t$ the exact value of $A_q(n, 2t; t)$ was known for more than forty years [1]. Via a computer search the cases $A_2(3k + 2, 6; 3)$ were settled in 2010 [5]. In 2015 the entire case $q = r = 2$ was resolved by continuing the original approach of Beutelspacher [11], i.e., by *considering* the set of holes in $(n - 2)$ -dimensional subspaces. Very recently, this was generalized to the consideration of the set of holes in $(n - j)$ -dimensional subspaces, where $j \leq t - 2$, and general q [12] so that we now know the exact values of $A_q(kt + r, 2t; t)$ in all cases where $t > \begin{bmatrix} r \\ 1 \end{bmatrix}_q := \frac{q^r - 1}{q - 1}$. Here, we streamline and generalize their approach leading to improved upper bounds on $A_q(n, 2t; t)$.

2. SUBSPACES WITH THE MINIMUM NUMBER OF HOLES

Definition 2.1. A vector space partition \mathcal{P} of \mathbb{F}_q^n has *hole-type* (t, s, m_1) , if it is of type $t^{m_t} \dots s^{m_s} 1^{m_1}$, for some integers $n > t \geq s \geq 2$, $m_i \in \mathbb{N}_0$ for $i \in \{1, s, \dots, t\}$, and \mathcal{P} is non-trivial.

Lemma 2.2. Let \mathcal{P} be a vector space partition of \mathbb{F}_q^n of hole-type (t, s, m_1) and $l, x \in \mathbb{N}_0$ with $\sum_{i=s}^t m_i = lq^s + x$. $\mathcal{P}_H = \{U \cap H : U \in \mathcal{P}\}$ is a vector space partition of type $t^{m'_t} \dots (s-1)^{m'_{s-1}} 1^{m'_1}$, for a hyperplane H with \widehat{m}_1 holes. We have $\widehat{m}_1 \equiv \frac{m_1 + x - 1}{q} \pmod{q^{s-1}}$. If $s > 2$, then \mathcal{P}_H is non-trivial and $m'_1 = \widehat{m}_1$.

PROOF. If $U \in \mathcal{P}$, then $\dim(U) - \dim(U \cap H) \in \{0, 1\}$ for an arbitrary hyperplane H . For $s > 2$, counting the 1-dimensional subspaces of \mathbb{F}_q^n and H , via \mathcal{P} and \mathcal{P}_H , yields

$$(lq^s + x) \cdot \begin{bmatrix} s \\ 1 \end{bmatrix}_q + aq^s + m_1 = \begin{bmatrix} n \\ 1 \end{bmatrix}_q \quad \text{and} \quad (lq^s + x) \cdot \begin{bmatrix} s-1 \\ 1 \end{bmatrix}_q + a'q^{s-1} + \widehat{m}_1 = \begin{bmatrix} n-1 \\ 1 \end{bmatrix}_q$$

for some $a, a' \in \mathbb{N}_0$. Since $1 + q \cdot \begin{bmatrix} n-1 \\ 1 \end{bmatrix}_q - \begin{bmatrix} n \\ 1 \end{bmatrix}_q = 0$ we conclude $1 + q\widehat{m}_1 - m_1 - x \equiv 0 \pmod{q^s}$. Thus, $\mathbb{Z} \ni \widehat{m}_1 \equiv \frac{m_1 + x - 1}{q} \pmod{q^{s-1}}$. For $s = 2$ we have

$$(lq^2 + x) \cdot (q + 1) + aq^2 + m_1 = \begin{bmatrix} n \\ 1 \end{bmatrix}_q \quad \text{and} \quad (lq^2 + x - m'_1 + \widehat{m}_1) \cdot (q + 1) + a'q^2 + m'_1 = \begin{bmatrix} n-1 \\ 1 \end{bmatrix}_q$$

leading to the same conclusion $\widehat{m}_1 \equiv \frac{m_1 + x - 1}{q} \pmod{q^{s-1}}$. \square

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¹The more general notation $A_q(n, 2t - 2w; t)$ denotes the maximum cardinality of a collection of t -dimensional subspaces, whose pairwise intersections have a dimension of at most w . Those objects are called *constant dimension codes*, see e.g. [6]. For known bounds, we refer to <http://subspacecodes.uni-bayreuth.de> [9] containing also the generalization to *subspace codes* of mixed dimension.

Lemma 2.3. *Let \mathcal{P} be a vector space partition of \mathbb{F}_q^n of hole-type (t, s, m_1) , $l, x \in \mathbb{N}_0$ with $\sum_{i=s}^t m_i = lq^s + x$, and $b, c \in \mathbb{Z}$ with $m_1 = bq^s + c$. If $x \geq 1$, then there exists a hyperplane \widehat{H} with $\widehat{m}_1 = \widehat{b}q^{s-1} + \widehat{c}$ holes, where $\widehat{c} := \frac{c+x-1}{q} \in \mathbb{Z}$ and $b > \widehat{b} \in \mathbb{Z}$.*

PROOF. Apply Lemma 2.2 and observe $m_1 \equiv c \pmod{q^s}$. Let the number of holes in \widehat{H} be minimal. Then,

$$\widehat{m}_1 \leq \text{average number of holes per hyperplane} = m_1 \cdot \frac{\binom{n-1}{1}_q}{\binom{n}{1}_q} < \frac{m_1}{q}. \quad (1)$$

Assuming $\widehat{b} \geq b$ yields $q\widehat{m}_1 \geq q \cdot (bq^{s-1} + \widehat{c}) = bq^s + c + x - 1 \geq m_1$, which contradicts Inequality (1). \square

Corollary 2.4. *Using the notation from Lemma 2.3, let \mathcal{P} be a non-trivial vector space partition with $x \geq 1$. For each $0 \leq j \leq s-1$ there exists an $(n-j)$ -dimensional subspace U containing \widehat{m}_1 holes with $\widehat{m}_1 \equiv \widehat{c} \pmod{q^{s-j}}$ and $\widehat{m}_1 \leq (b-j) \cdot q^{s-j} + \widehat{c}$, where $\widehat{c} = \frac{c + \binom{[1]_q \cdot (x-1)}{j}}{q^j}$.*

Lemma 2.5. *Let \mathcal{P} be a vector space partition of \mathbb{F}_q^n with $c \geq 1$ holes and a_i denote the number of hyperplanes containing i holes. Then, $\sum_{i=1}^c a_i = \binom{n}{1}_q$, $\sum_{i=1}^c ia_i = c \cdot \binom{n-1}{1}_q$ and $\sum_{i=2}^c i(i-1)a_i = c(c-1) \cdot \binom{n-2}{1}_q$.*

PROOF. Double-count the incidences of the tuples (H) , (B_1, H) , and $(\{B_1, B_2\}, H)$, where H is a hyperplane and $B_1 \neq B_2$ are points contained in H . \square

Lemma 2.6. *Let $\Delta = q^{s-1}$, $m \in \mathbb{Z}$, and \mathcal{P} be a vector space partition of \mathbb{F}_q^n of hole-type (t, s, c) . Then, $\tau_q(c, \Delta, m) \cdot \frac{q^{n-2}}{\Delta^2} - m(m-1) \geq 0$, where*

$$\tau_q(c, \Delta, m) = m(m-1)\Delta^2q^2 - c(2m-1)(q-1)\Delta q + c(q-1)(c(q-1) + 1).$$

PROOF. Consider the three equations from Lemma 2.5. $(c - m\Delta)(c - (m-1)\Delta)$ times the first minus $(2c - (2m-1)\Delta - 1)$ times the second plus the third equation, and then divided by $\Delta^2/(q-1)$, gives

$$(q-1) \cdot \sum_{h=0}^{\lfloor c/\Delta \rfloor} (m-h)(m-h-1)a_{c-h\Delta} = \tau_q(c, \Delta, m) \cdot \frac{q^{n-2}}{\Delta^2} - m(m-1)$$

due to Lemma 2.2. Finally, we observe $a_i \geq 0$ and $(m-h)(m-h-1) \geq 0$ for all $m, h \in \mathbb{Z}$. \square

Lemma 2.7. *For integers $n > t \geq s \geq 2$ and $1 \leq i \leq s-1$, there exists no vector space partition \mathcal{P} of \mathbb{F}_q^n of hole-type (t, s, c) , where $c = i \cdot q^s - \binom{s}{1}_q + s - 1$.*

PROOF. Assume the contrary and apply Lemma 2.6 with $m = i(q-1)$. Setting $i = s-1-y$ we compute

$$\tau_q(c, \Delta, m) = -q\Delta(y(q-1) + 2) + (s-1)^2q^2 - q(s-1)(2s-5) + (s-2)(s-3).$$

Using $y \geq 0$ we obtain $\tau_2(c, \Delta, m) \leq s^2 + s - 2^{s+1} < 0$. For $s = 2$, we have $\tau_q(c, \Delta, m) = -q^2 + q < 0$ and for $q, s > 2$ we have $\tau_q(c, \Delta, m) \leq -2q^s + (s-1)^2q^2 < 0$. Thus, Lemma 2.6 yields a contradiction. \square

Theorem 2.8. *For integers $r \geq 1$, $k \geq 2$, and $z, u \geq 0$ with $t = \binom{r}{1}_q + 1 - z + u > r$ we have $A_q(n, 2t; t) \leq lq^t + 1 + z(q-1)$, where $l = \frac{q^{n-t} - q^r}{q^t - 1}$ and $n = kt + r$.*

PROOF. Assume the existence of a non-trivial vector space partition \mathcal{P} of type $t^{m_t} 1^{m_1}$ of \mathbb{F}_q^n with $m_t = lq^t + x$, where $x = 2 + z(q-1)$. Since $m_t \cdot \binom{t}{1}_q + m_1 = \binom{n}{1}_q$, we have $m_1 = bq^t + c$, where $b = \binom{r}{1}_q$ and $c = -\binom{t}{1}_q(x-1)$. Apply Corollary 2.4 with $s = t$ and $j = t - z - 1$ on \mathcal{P} . The $(n - t + z + 1)$ -dimensional subspace U contains $L \leq (2z - u)q^{z+1} + \frac{-\binom{t}{1}_q(x-1) + \binom{t-z-1}{1}_q(x-1)}{q^{t-z-1}} \leq zq^{z+1} - \binom{z+1}{1}_q + z$ holes. For $z = 0$ this number is negative and for $z \geq 1$, we can apply Lemma 2.7 using $L \equiv zq^{z+1} - \binom{z+1}{1}_q + z \pmod{q^{z+1}}$ (see Lemma 2.2). \square

The known constructions for partial t -spreads give $A_q(kt + r, 2t; t) \geq lq^t + 1$, see e.g. [1] (or [11] for an interpretation using the more general multilevel construction for subspace codes). Thus, Theorem 2.8 is tight for $t \geq \binom{r}{1}_q + 1$, c.f. [12, Lemma 9].

Theorem 2.9. For integers $r \geq 1$, $k \geq 2$, and $z \geq 0$ with $t = \begin{bmatrix} r \\ 1 \end{bmatrix}_q + 1 - z > r$, $n = kt + r$, and $l = \frac{q^{n-t} - q^r}{q^t - 1}$, we have $A_q(n, 2t; t) \leq lq^t + q^{r+1} - \left\lfloor \frac{1}{2} + \frac{1}{2} \cdot \sqrt{4q^{r+1}(q^{r+1} - (z+r)(q-1) - 1) + 1} \right\rfloor$.

PROOF. Assume the existence of a non-trivial vector space partition \mathcal{P} of type $t^{m_t} 1^{m_1}$ of \mathbb{F}_q^n with $m_t = lq^t + x$, where $x \geq 1$. Since $m_t \cdot \begin{bmatrix} t \\ 1 \end{bmatrix}_q + m_1 = \begin{bmatrix} n \\ 1 \end{bmatrix}_q$, we have $m_1 = bq^t + c$, where $b = \begin{bmatrix} r \\ 1 \end{bmatrix}_q$ and $c = -\begin{bmatrix} t \\ 1 \end{bmatrix}_q(x-1)$. Apply Corollary 2.4 with $s = t$ and $j = t - r - 1$ on \mathcal{P} . The $(n - t + r + 1)$ -dimensional subspace U contains $L \leq (z+r)q^{r+1} + \frac{-\begin{bmatrix} t \\ 1 \end{bmatrix}_q(x-1) + \begin{bmatrix} t-r-1 \\ 1 \end{bmatrix}_q(x-1)}{q^{t-r-1}} = (z+r)q^{r+1} - \begin{bmatrix} r+1 \\ 1 \end{bmatrix}_q(x-1)$ holes. Due to Lemma 2.2 we have $L \equiv -\begin{bmatrix} r+1 \\ 1 \end{bmatrix}_q(x-1) \pmod{q^{r+1}}$. Next, we will show that $\tau_q(c, \Delta, m) \leq 0$, where $\Delta = q^r$ and $c = iq^{r+1} - \begin{bmatrix} r+1 \\ 1 \end{bmatrix}_q(x-1)$ with $1 \leq i \leq z+r$, for a suitable x and $m \geq 1$. Applying Lemma 2.6 then gives the desired contradiction, so that $A_q(n, 2t; t) \leq lq^t + x - 1$.

We choose² $m = i(q-1) - x + 2$. With this, solving the quadratic equation $\tau_q(c, \Delta, m) = 0$ for x gives $x_0 = \bar{\Delta} + \frac{1}{2} \pm \frac{1}{2}\theta(i)$, where $\bar{\Delta} = q\Delta = q^{r+1}$ and $\theta(i) = \sqrt{4\bar{\Delta} \cdot (\bar{\Delta} - i(q-1) - 1) + 1}$. Since $\lim_{x \rightarrow \infty} \tau_q(c, \Delta, m) = \infty$, we have $\tau_q(c, \Delta, m) \leq 0$ for $|2x - 2\bar{\Delta} - 1| \leq \theta(i)$. We need to find an integer x such that this inequality is satisfied for all $1 \leq i \leq z+r$. The strongest restriction is attained for $i = z+r$. Since $z \leq \begin{bmatrix} r \\ 1 \end{bmatrix}_q - r$, we have $\theta(i) \geq \theta(z+r) \geq 1$, so that $\tau_q(c, \Delta, m) \leq 0$ for $x = \bar{\Delta} - \lfloor -\frac{1}{2} + \frac{1}{2}\theta(z+r) \rfloor$. With respect to Lemma 2.6 we remark that $-m(m-1) < 0$ for all $m \in \mathbb{Z} \setminus \{0, 1\}$. So, it remains to verify $\tau_q(c, \Delta, m) < 0$ for $m \in \{0, 1\}$. If $i < z+r$ this is true due to $\theta(i) > \theta(z+r)$, so that we assume $i = z+r$. Due to Theorem 2.8 it suffices to consider the cases $x \leq 1 + z(q-1)$. Thus $m \geq r(q-1) + 1 \geq 2$. \square

For the special case $t = r + 1$, Theorem 2.9 is equivalent to [4, Corollary 8], which is based on [3, Theorem 1B]. And indeed, our analysis is very similar to the technique³ used in [3]. The new ingredients essentially are lemmas 2.2 and 2.3. Postponing the details and proofs to a more extensive and technical paper, we state:

- $2^4l + 1 \leq A_2(4k + 3, 8; 4) \leq 2^4l + 4$, where $l = \frac{2^{4k-1} - 2^3}{2^4 - 1}$ and $k \geq 2$, e.g., $A_2(11, 8; 4) \leq 132$;
- $2^6l + 1 \leq A_2(6k + 4, 12; 6) \leq 2^6l + 8$, where $l = \frac{2^{6k-2} - 2^4}{2^6 - 1}$ and $k \geq 2$, e.g., $A_2(16, 12; 6) \leq 1032$;
- $2^6l + 1 \leq A_2(6k + 5, 12; 6) \leq 2^6l + 18$, where $l = \frac{2^{6k-1} - 2^5}{2^6 - 1}$ and $k \geq 2$, e.g., $A_2(17, 12; 6) \leq 2066$;
- $2^7l + 1 \leq A_2(7k + 5, 14; 7) \leq 2^7l + 17$, where $l = \frac{2^{7k-2} - 2^5}{2^7 - 1}$ and $k \geq 2$, e.g., $A_2(19, 14; 7) \leq 4113$;

c.f. the web-page mentioned in footnote 1 for more numerical values.

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²The choice for m is obtained by solving $\frac{\partial \tau_q(c, \Delta, m)}{\partial m} = 0$, i.e., we minimize $\tau_q(c, \Delta, m)$, and up-rounding the unique solution.

³Actually, their analysis grounds on [13] and is strongly related to the classical second-order Bonferroni Inequality [2, 7, 8] in Probability Theory, see e.g. [10, Section 2.5] for another application for bounds on subspace codes.